# An Industrial Case Study of Coman's Automated Task Detection Algorithm: What Worked, What Didn't, and Why

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Abstract—Programmers need explicit tool support for software maintenance tasks, and a prerequisite for this is an understanding of where the boundaries between distinct tasks lie. Asking developers to indicate manually when they switch tasks is disruptive to their normal work flow, so researchers have sought ways to infer task boundaries automatically based on the content of the interaction histories with the IDE. Coman previously reported a fully automated algorithm that achieved 80% accuracy in a lab validation study. In this paper, we evaluate the use of this algorithm within an industrial setting. We found two problems: first, a large number of the tasks identified are in fact only sessions or subparts of a larger task; second, the demonstrable effects of interruptions are not considered. We argue that the problem of task boundary detection consists of two sub-problems: first, detecting task sessions; and second, linking task sessions. Coman's algorithm only partially addresses the first, and ignores the second.

Keywords-task boundary detection; case study;

### I. Introduction

Software maintenance tasks, such as fixing bugs and adding features, are typically highly labour intensive. Explicit tool support for tasks is often weak, and different kinds of subtasks require a variety of skills, artifacts, and tools [1], [2], [3]. Moreover, a task may be interrupted or switched out, causing extra effort in managing task context [4], [5].

Recently, practitioners and researchers have begun exploring ways to provide explicit support for typical maintenance tasks. For example, researchers have proposed to improve IDEs by providing task-friendly views, to recommend task relevant information [3] and to build task-aware software development [6]. A fundamental problem for these task-aware applications is to detect task boundaries, that is when a task is started, interrupted, resumed, or finished. Without an understanding of where the boundaries lie, supporting tools may unknowingly mix task contexts together and provide inappropriate information at inappropriate time.

The problem of task boundary detection is challenging [6]. The ways to start a task vary from the task type, the comprehension strategy, and programmers expertise. Also,

various cases of interruption and task switching make the problem more complicated.

To our knowledge, Coman et al. have proposed the only automated algorithm for detecting task boundaries from interaction histories [7], [8]. In a lab validation study, this algorithm has shown 80% success in identifying the number of tasks. However, industrial software development is quite different from lab setting. Software systems are often much larger, tasks are more complex, and interruptions are common. In this work, we have sought to evaluate how Coman's algorithm may perform within an industrial setting.

We performed an industrial case study involving six professional software developers working on their normal real-world tasks. From over a month's worth of raw data, we examined 12 randomly chosen person-days in detail. We applied Coman's algorithm on the interaction histories captured by the IDE, and compared the output against the "ground truth" as reported by the developers themselves. When we noticed surprising results, we further examined the details of the data to obtain insights into the algorithm and the problem space, and considered possible ways to improve the algorithm.

We first discuss the problem in Section II. We next review Coman's algorithm and lab study, and present our research questions in Section III. We present the case study design in Section IV and the results in Section V. We discuss related work in Section VI, and threats to validity in Section VII. Finally, we conclude in Section VIII.

# II. THE PROBLEM AND CHALLENGE

A programmer's work is driven by tasks, such as fixing bugs or adding new features. In a diary study of 13 industrial developers, 13 types of tasks are found, such as estimation/investigation, code and high-level test [9]. However, it should be pointed out that there is no generally accepted categorization of developer task types. In our work here, we consider only software maintenance tasks that include the modification of code. Such tasks are usually categorized into four kinds: adaptive, corrective, perfective, and preventative.

Programmers may spend a lot of effort solving a given software maintenance task. They need to go through several

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different stages, such as coming to an understanding of desired behavior, gathering relevant information, and testing hypotheses. At each of these stages, they usually need to develop goals, hypotheses, and ask questions [1], [2]. Moreover, in a large software system, task-related information is often scattered around in different places [3]. Developers may need to navigate and search a lot to find task relevant artifacts. Tasks may be decomposed into subtasks. For example, when a current task requires code change by another programmer, then a subtask may be created for the other programmer [5].

A task may also be interrupted or switched out, causing extra effort in managing task context. Programmers may interrupt themselves or be interrupted [4]. If the current task changes to a lower priority, it may be switched out in favor of a higher-priority task. Both in cases of interruption and task switching, the context of tasks need to be saved and loaded, resulting in much mental effort. Task switching is considered by programmers as a serious problem in their work [5].

To help programmers, researchers have tried to provide explicit task support for software development. Murphy et al. argued that task structure can be used to improve IDEs by providing task-friendly views, to recommend task relevant information, and to build a group memory for collaboration [3]. Robillard et al. proposed to support task-aware software development environment based on navigation analysis [6].

However, the problem of task boundary detection remains challenging. There are many ways a given developer might choose to start a given task, making it hard to detect the starting point. Moreover, there are various reasons why a task might be interrupted or switched out, making it hard to identify the transition points between tasks.

Comprehension strategy, task type, programmer expertise, and other factors may all affect how a task is started. Some developers may start a task by building and exploring hypotheses (top-down) while some may by reading code (bottom-up) [1]. Some may begin with searching while some with navigation [10], [11]. For corrective tasks, the common first step may be recreating the problem; for perfective tasks, the first step may be comprehending the desired behavior. Even for the same initial step, the methods used may be different: for example, to comprehend the desired behavior, some developers may use debugging while some may prefer searching the documentation.

Also, programmer expertise affects how a task may be started. An experienced programmer with a good understanding of the code may go directly into the code locations that need to be changed, while a newcomer may have to explore code for hours just to find an appropriate starting point.

Various cases of interruption and task switching also make the problem hard to solve. Many events may trigger the transition of tasks: changing of task priority, getting blocked, getting a request from colleagues, lunch time/conference, or simply getting tired of the current task. These events are often not explicitly recorded anywhere, therefore providing little information that a task boundary detection algorithm may rely on.

### III. COMAN'S ALGORITHM AND LAB STUDY

Coman et al. proposed an algorithm that detects task boundaries automatically based on programmer interaction histories. The main ideas behind the algorithm are:

- Each task has a set of artifacts that are essential for performing the task, namely the *task core*.
- When programmers work on a task, there is a time period when the task core is accessed intensively at approximately the same time.

Based on these assumptions, the algorithm first computes the *time intervals of intensive access* (TIIAs) for each method — these are the time periods with intensive access to the method. Time moments with a large number of TIIAs are then identified as task core moments. Using these task core moment as seeds, the TIIAs are grouped into task subsections based on the temporal distances. The resulting task subsections are finally identified as the tasks.

Since we refer to the details of the algorithm in later discussions, we now describe the main steps and parameters of Coman's algorithm:

# 1) Compute TIIA(m)

Based on interaction histories, the algorithm computes TIIAs of each method based on degree of access (DOA), which is defined as  $DOA(m,t) = \frac{AT(m,t)}{t-t_0}$ , where AT(m,t) is the amount of time that a method is accessed and  $(t-t_0)$  is the total interval of time period. By definition, the value of DOA is between 0 and 1. A TIIA of a method starts when it is first accessed, and ends when its DOA decreases below a stated threshold, th. The higher the th, the stricter for being considered as "intensive access", therefore the more likely to end an existing TIIA in case of no access. th is a parameter for tuning.

### 2) Form TIIA time series

Once all of the TIIAs have been computed for all of the methods, the number of TIIAs at each time can be computed. This forms a TIIA time series.

## 3) Smooth TIIA time series

The TIIA time series is smoothed using weighted central moving average (WCMA). Every data point is computed as an average of the current point and a number of data points on either side, with weights decreasing as the temporal distance increases. Smoothing transforms the TIIA time series into a continuous line so that peaks and valleys can be identified later on.

The number of data points used in WCMA is a parameter for tuning, namely ta.

# 4) Identify task core moments

Peaks on the TIIA time series indicate moments when a large number of methods are accessed within a short period. By assumption, these methods are task cores and the peaks are task core moments.

Only peaks with height greater than a threshold tp are considered as task core moments. The height of a peak is computed as the difference between the peak and the

higher of the two adjacent valleys. tp is a parameter for tuning.

# 5) Expand task core moments to subsections The task core moments are then expanded by adding all of the TIIAs one by one. Each TIIA is grouped with the task subsection that has the smallest distance to it. The distance of a TIIA and a task subsection is defined as the spanning difference of the task subsection if adding the TIIA.

Each task subsection is then identified as a task.

The algorithm has three parameters that can be tuned: th for the DOA threshold, ta for the number of data points (or the time) on each side in smoothing, and tp for peak height threshold. Coman et al. state that parameter tuning is important for applying the algorithm, but they do not provide specific guidelines for doing so.

To investigate the proposed algorithm empirically, Coman et al. performed a laboratory experiment. Three students were allowed 70 minutes to solve five maintenance tasks for the Paint program, which has 9 classes and 503LOC. Both the tasks and the Paint code were developed by other researchers. In this study, the parameter values used were  $th=\frac{2}{3}*median(accesses)*\frac{1}{10},\ ta=250(4min),\$ and th=0.2, where median(accesses) is the median length of the accesses. The definition of th takes the characteristics of programmer accesses into consideration. If a programmer tends to have long accesses, a single lack of access would not make th decreases rapidly and end the current TIIA. Parameter ta was set to 250, which equals roughly 4 minutes on either side. Peak threshold th was set to 0.2.

In this lab study, the algorithm was able to correctly identify 9 out of the 11 tasks that were attempted, achieving a success rate of 81%. All of the task subsections that were identified as tasks did indeed correspond to a "ground truth" task. The error of 19% came from that two small tasks were not identified by the algorithm.

# A. Research questions

We know from experience that industrial software development is often quite different from the artificial setting in the lab. In industry, the software systems are often very large, the tasks to be solved are complex, and interruptions are common. All these factors may cause techniques that work well in the lab to behave differently in an industrial setting.

With this in mind, we chose to investigate two research questions:

- Q1: How does Coman's algorithm perform in an industry setting?
- Q2: If the algorithm performs differently than expected, why?

# IV. CASE STUDY DESIGN

To answer our research questions, we performed an industrial case study of detecting task boundaries. The study was

TABLE I BACKGROUND OF PARTICIPANTS

Prog	Proj	Years	Years	Years	
		in Proj	in Java	in Eclipse	
N1	H3.DBM	0.8	4	2	
P1	H3.DBM	0.8	4	2	
P2	H3.DBM	0.8	5	2	
P3	H3.SYS	2	5	2	
P4	H3.KNW	0.6	6	2	
P5	Learning	0	1	2	
P6	Learning	0	2	2	

performed in Heweisoft, a software company located in Shanghai, China whose main expertise is in building information system for enterprise and government.

Six programmers from the R&D department participated in this study; we shall refer to them as P1, P2, etc. All of them were actively involved in the development of an internal platform called H3, a middleware system for distributed enterprise applications. The programmers were working on different sub-projects of H3: P1 and P2 were developing H3.DBM, a distributed database engine of about 300KLOC, P3 worked for H3.SYS, a system management component of about 50KLOC, and P4 was in H3.KNW subproject, a knowledge management system of about 100KLOC. At the time of the study, the H3 platform was being used by two other projects for building business applications and was in active maintenance.

Before the study, each programmer completed a questionnaire about their background and experience. The results are shown in Table I. At the time the study was performed, P1, P2, P3, and P4 had been on the H3 team for about two years, and all were considered to have good understanding of the whole project. P1, P2, and P4 had switched to their current sub-project six to eight months previously; all had become experienced developers of the sub-project. P5 and P6 were newcomers to both the project and the team, having just joined the company the previous week. During the first period of learning, they were assigned some tasks to become familiar with the software system. To ensure that the level of granularity of a task in our study is comparable to Coman's study, we pointed out to programmers that typical software maintenance tasks include adding a new feature and fixing a bug.

For each programmer, we captured their interaction histories for a month using a special-purpose plug-in tool we created for their development environment, Eclipse. Captured events have the format (time, action, artifact) which indicate respectively the time at which a developer acted on an artifact (method or file), whether the action was a view or a change, and the identity of the artifact. The tool required no input from the user and was invisible to the programmers, although of course they were all aware of its existence and use.

For each programmer, we randomly picked three days before the start of the study as the days for applying the task splitting algorithm. The developers did not know ahead of the time which days had been picked. At the end of each

TABLE II BASIC RESULTS

	#Tasks self-			# Interaction Events		
	reported/predicted			on the day		
	("-": no report)			-		
Prog	D1	D2	D3	D1	D2	D3
P1	1/16	2/23	-	7141	22221	-
P2	2/15	3/15	-	10485	10072	-
P3	3/0	2/1	-	1701	1367	-
P4	2/2	-	-	752	-	-
P5	1/1	1/12	1/4	478	6444	1060
P6	-	1/22	2/5	-	8527	1496

of those days, developers were asked to send email detailing what tasks they had performed that day, and how their time had been broken up into tasks. The email asked two questions:

1) What tasks have you performed today? 2) How was your time segmented into these tasks? The developers were asked to do this at the end of these chosen days so that the knowledge of what they had done that day would still be fresh in their memories. In case that programmers did not response, which we may detect in the next morning, we did not ask developers to recall yesterday's tasks again, since we thought that it might be too hard on them. So we ignored all the non-reponses.

For each of the selected days, we applied the algorithm and compared the results with the number of self-reported tasks, which served as our ground truth. We first ran the algorithm with the same parameter values had been used in the lab study; we refer to these as the *baseline values* for the parameters. We then tuned the parameters to find values that best fit our data. Since Coman et al. provided no guidance on parameter tuning, we did so based on our own understanding of the algorithm: for each combination of parameter values, we applied the algorithm and chose the parameter values that best fit with the self-reported number.

We did not verify the results of Coman's algorithm with programmers since we thought that might affect programmers' response. Verification for this study is better done when programmers still have fresh memory about their work, which means at the end of all the selected days. However, showing the results on multiple days may make programmers speculate the purpose of the study and tend to report numbers that are close to that.

### V. CASE STUDY RESULTS

We received a total of 12 email messages from the developers describing their tasks and time spent, although 18 had been expected. Table II shows these results, as well as the number of interaction events on those days to give a rough idea of the activity level.

We can see from the table that the number of tasks everyday is small, ranging from one to three. Most programmers reported only a very rough time segmentation, often one in the morning, and another in the afternoon. We consider this understandable since it is much more difficult to recall the details of time periods than the number of task. Most tasks reported were "debugging and fixing X bug", "deploying X to server", "learning X module", and "preparing training system", so we consider the granularity of tasks in our study is quite similar to Coman's study. For P5 and P6, since they were newcomers, the tasks they reported were mostly "learning X".

A. Q1: How does Coman's algorithm work in an industry setting?

To answer Q1, we applied the algorithm and compare the results with self-reported data, the ground truth. We ran the algorithm under two situations, with baseline parameters and with parameter tuning.

1) Results with baseline parameters: For each of the 12 days, we apply the task splitting algorithm using the baseline parameters:  $th = \frac{2}{3} * median(accesses) * \frac{1}{10}$ , ta = 250, and tp = 0.2. Table II shows the results.

We can see that the number of tasks detected fits poorly with the self-reported number overall. The total number of tasks with the baseline parameters was 112, while the total number of the self-reported value was only 21.

The algorithm's accuracy varied over the different days of the study. On two days, P4.D1 and P5.D1, the number of tasks detected was the same as the self-reported. But in half of the 12 total cases, such as P1.D1 and P1.D2, the number of tasks found by the algorithm was much larger than the self-reported number. In the case of P3.D1, no tasks were detected. Further analysis shows that this was because no peak has height greater than the threshold value (the highest peak has height of 0.14).

- 2) Results of parameter tuning: We tried to tune the parameters to better fit our data. Since the baseline parameter values worked well in the lab study, we judged it likely that the best parameters for our study would be somewhere close by. Consequently, we performed the parameter tuning around baseline values. For each parameter, we pick a few values below and a few above the baseline values. In more detail, we tune the parameters th, ta, and tp as follows:
  - DOA threshold: th

th is the threshold for the level of DOA that ends a TIIA. It is set to  $th = \frac{2}{3}*median(accesses)*\frac{1}{10}$  in the baseline, where median(accesses) is the median duration between accesses. A larger value for th means shorter TIIAs, and thus fewer tasks.

In our case study, the median(accesses) of the 12 days have total four values, four of 1, four of 2, one of 3, and one of 6. This means the range of th is from 0.067(th=1) to 0.4 (th=6). So following the rationale of "close to the baseline", we choose the following values:

 $th = 0.04, \frac{2}{3} * median(accesses) * \frac{1}{10}, 0.2, 0.4.$ 

• Smoothing parameter: ta

ta defines the number of data points on either side involved in smoothing. The larger the ta is, the smoother the time series will be, and thus the peak height will be smaller and fewer tasks will be identified.

The baseline value of ta is 250, roughly 4 minutes on either side, so we choose following ta values:

ta = 180, 250, 300, 420(3min, 4min, 5min, 7min)

• Peak height threshold: tp

tp is the threshold for peak identification in the time series. The larger tp is, the harder it is for a peak to be considered as a task core moment, thus leading to fewer peaks and fewer tasks.

The baseline value of tp is 0.2, so we try following values:

$$tp = 0.1, 0.2, 0.4, 0.6$$

For each combination of parameter values, we ran the task splitting algorithm. The results are shown in Figure 1: the 12 series represent the 12 programmer days, the X-axis represents particular combinations of parameter value, and the Y-axis represents the number of tasks being detected. There are a total of 64 data points on the X-axis, since each of the three parameters has four values. These data points are sorted in the order of th, ta, and tp, starting from th = 0.04, ta = 180(3min), tp = 0.1 and ending at th = 0.4, ta = 420(7min), tp = 0.6.

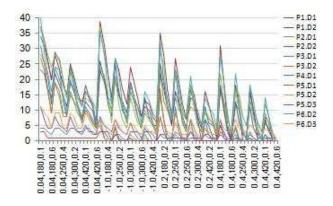


Fig. 1. Parameter tuning

As we can see from the figure, some series are quite spiky, such as P1.D1 and P6.D2, while some are smoother, such as P5.D1. We notice a common pattern for a large number of series: a total of 16 spikes on a series, and every 4 form a group; within each group, the height of the four spikes decrease and each spike is sharp; and across the groups, the overall height decreases slightly. This pattern suggests that the effect of parameters has some regularity.

The rapidly decreasing height of each spike indicates that the effect of tp on the algorithm is strong. Increasing tp lifts the threshold of being identified as peaks, and greatly reduces the number of tasks greatly.

Within each group, the decreasing height of the four spikes shows the effect of ta, since the only difference between the four spikes is ta. The bigger ta, the stronger the smoothing, which causes lower peaks and higher valleys. The rapidly lowering spikes within a group suggests that smoothing has a strong effect on the algorithm.

Across the four groups, the change of the overall height is the result of changing th, as the only difference between the groups is th, the threshold of DOA. We found that the four groups look quite similar, with only a slight decrease of the overall height. This suggests that the effect of th on the algorithm is relatively small,

For each parameter combination, we compare the number of tasks being detected with the self-reported number. The parameter sets that have the smallest total difference for all the 12 cases are chosen as the best fits.

We found that the parameters that fit best are (th=0.4,ta=180,tp=0.6)(51), and (th=0.4,ta=420,tp=0.4)(62). In both cases, the total difference of number of tasks identified is 16. The total number of self-reported tasks is 21, so the overall error rate is about 16/21=76%. In both parameter sets, no tasks are detected in half of the 12 programmer days, or that the parameters are over-tuned for these cases. This suggests that parameter tuning for the algorithm may not be as simple as finding a universal set of values. It may depend on factors such as programmer, task, and the software system.

We feel it is important to pay attention to the underlying meaning of the parameter values while tuning the parameters. For example, smoothing means that information is lost. Setting ta=600 means 10 minutes on each side participated in smoothing, which means 20 minutes in total. However, 20 minutes may be long enough to solve a small task. So would such small tasks be washed away in such smoothing? Also for tp, increasing its value means that higher peaks are considered as task cores, which also means that more methods should be accessed intensively in a shorter time period to be identified as a task core. It is unclear that this is a reasonable assumption.

3) Answer to Q1: From the results presented above, we can see that the performance of the algorithm in our industrial setting is much different from that observed in the lab setting in the original study. With the same parameters as in the lab study, the total number of tasks being detected was about 5 times that of the self-reported number. Results vary with cases: some are the exactly same, some are below, but most are well above the self-reported number. After tuning the parameters, the best cases still have error rate of about 76%. The parameter tuning may be more complicated than finding a universal applicable parameter set.

B. Q2: If the algorithm performs differently than expected, why?

Since the results of the Coman algorithm were quite different in our industrial study than the original lab results, we sought to understand why this might be the case by performing a deeper analysis of the data and the assumptions of the algorithm.

1) Large number of predicted tasks: We wondered why, in many cases, Coman's algorithm tended to detect many more tasks than were reported by the developers themselves. We examined the details of the data: the TIIA time series, the peaks, the task core, and task subsections. We used P1.D1 under the baseline parameter as our example. On that day programmer P1 reported performing only one task.

Figure 2 shows the TIIA time series of P1.D1 before and after smoothing. 16 peaks on the smoothing series are identified as the task core moments, yet no single peak is dominant.

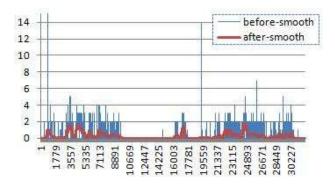


Fig. 2. TIIA time series being smoothed

We further looked into the task cores, as shown in Figure 3. In this figure, each method that is ever included in at least one task core is plotted on the X-axis. Each row represents a task as detected from the algorithm. A dot at (X, Y) represents that method X is included in the core of task Y. As we can see, the size of the task core is small, about 2 methods on average. There are small overlaps between task cores: *e.g.*, the task core of S7 is exactly the same as S14, and is a subset of the task core of S8.

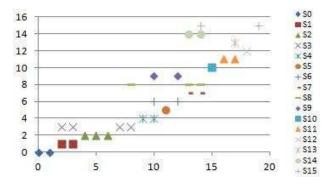


Fig. 3. Methods in task cores

Figure 4 shows the artifacts included in each task that was detected. This figure is similar to Figure 3 except that X-axis is files and a dot at (X,Y) means that file X is included in the task of Y. We plot files instead of methods since the number of methods is quite large. We can see from the figure that artifacts between tasks overlap greatly. Many consecutive tasks have a large number of common artifacts. For example, all of the artifacts in S9 are contained in S10, same is for S13 and S14.

We used the same methods examining the other days that have a large number of tasks. We observed that the data shows a common pattern: no dominant peak, small size of task cores, small overlap between task cores, and large overlap between task artifacts. Such pattern suggests that a task is completed through multiple stages (multiple peaks), with a set of methods that are essential to the stage being accessed intensively approximately the same time (peak); the methods core to each stage may vary, but since for solving one task, the sets of artifacts involved in each stage have large overlaps.

We wonder what these stages mean. Do they correspond to

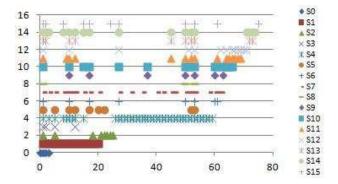


Fig. 4. Files in tasks

the typical stages of solving a complex task? For example, when fixing a bug, a programmer may first analyze the behaviour, then change the code, and finally check other code to make sure no new bugs are introduced. Can the first stage as detected by the algorithm be "analyzing behaviour"? More studies into this are needed.

So our data shows that the underlying assumption that artifacts that are accessed approximately the same time are a task core can be problematic in industrial software development. For complex tasks and under normal time pressure, a task is often completed through multiple stages, and artifacts essential to each stage are intensively accessed at the same time. In the lab study, since the software system is small and the task is simple, plus programmer working under very strict time pressure, the multiple stages are combined.

2) The effect of interruptions: Interruptions to the work flow of a task are not explicitly considered by Coman's algorithm; yet, it seems clear that interruptions can and do have important bearing on real-world task performance.

In the algorithm, a period of "no access" for an method will decrease the DOA and may end the TIIA. When an interruption occurs, the DOA of all the artifacts decreases rapidly, which may result the ending of a large number of TIIAs, forming a gap in the TIIA time series. The gap may result in a valley in the smoothed time series. The more valleys, the more possible peaks, and finally the more tasks. Therefore, interruptions may affect the number of predicted tasks.

Also in the algorithm, the task boundaries are determined by the overlap of TIIAs. Interruptions may cause large gaps of TIIAs. According to the algorithm, TIIAs before the gap will be grouped into some task before the gap, even though the core of the real task is after the gap. So, interruptions may also affect boundaries of the task.

To assess the effect of interruptions on the number of tasks, we decided to artificially "shrink" the interruptions in the interaction history and compare the number of tasks before and after the shrinking. For example for P1.D1, if we shrink all the interruption longer than five minutes to one minute, then the number of tasks is reduced from 16 to 13. Table III shows the difference. We can see that the effect of interruptions does indeed affect the output of the algorithm: in seven out of the

TABLE III #TASKS CHANGE WHEN INTR SHRINK

	#Tasks before → #Tasks after			
Prog	D1	D2	D3	
P1	$16 \rightarrow 13$	$23 \rightarrow 22$	-	
P2	15 → 14	$15 \rightarrow 15$	-	
P3	$0 \rightarrow 0$	$1 \rightarrow 0$	-	
P4	$2 \rightarrow 1$	-	-	
P5	$1 \rightarrow 1$	$12 \rightarrow 12$	$4 \rightarrow 3$	
P6	-	$22 \rightarrow 20$	$5 \rightarrow 4$	

TABLE IV
TASK BOUNDARIES OVERLAP WITH INTR

	#Tasks overlap / #Tasks			
Prog	D1	D2	D3	
P1	8/16	6/23	-	
P2	5/15	6/15	-	
P3	<b>-/</b> 0	<b>-</b> /1	-	
P4	1/2	-	-	
P5	-/1	9/12	2/4	
P6	-	10/22	4/5	

12 cases, the number of tasks decreased.

To see the effect of interruptions on the task boundary, we compared the boundaries of tasks with the interruptions within the interaction history. For example for P1.D1, for the total 16 tasks as detected, 8 of them have a starting time that overlaps with one of the top 20 interruptions.

Table IV shows the number of overlaps between the task boundary and interruptions for all the 12 days. Overall, one third to half of the task boundaries overlap with the top 20 interruptions. In some cases, such as P5.D2 and P6.D3, the degree of overlapping is large.

We also note that when the parameters are changed, overlapping is still present. For example, Table V shows the number of overlaps for P1.D1 under different parameters. We can see that about half of the starting time of the tasks overlap with the interruptions.

TABLE V INTR OVERLAP WHEN PARAMS CHANGE

	#Tasks overlap / #Tasks			
P1.D1	ta = 180	ta = 250	ta = 300	ta = 420
tp = 0.2	8/20	8/16	9/14	5/7
tp = 0.4	6/12	6/10	6/8	3/5
tp = 0.6	5/8	4/7	3/4	3/5

- 3) Answer to Q2: From the above discussions, we can see that there are two problems with Coman's algorithm that may account for its inconsistent performance between an industrial setting and a lab setting:
- 1) The assumption that the task core is accessed intensively at approximately at the same time is unrealistic. The tasks identified under such assumption may be just stages of solving a complex task. Due to the artificial lab setting, the stages of solving a simple task are clustered. But in an industrial setting,

this assumption does not hold and resulted in many more tasks being identified.

2) The effects of interruptions on the algorithm were ignored, yet interruptions are common in real-world development and affect the performance of the algorithm. The lab setting assumed no interruptions, but in an industrial setting with many interruptions, the effects were significant.

### C. Discussion

1) Insights into the problem: Interruptions and task switching are a fact of life for software developers, and result in fragmented task sessions. Therefore we consider that the problem of identifying the task boundaries in fact consists of two sub-problems. a) detecting task sessions — time periods when programmer is working on a single task continuously, and b) linking related task sessions together, such as grouping task sessions that belong to a same task. Obviously, solving the first sub-problem is the base for solving the second one. But if the second sub-problem is not solved, task sessions of the same task may be incorrectly identified as two different tasks, therefore leaving the larger problem unsolved.

The algorithm proposed by Coman et al. partially addresses the first sub-problem. The key of detecting task sessions is to identify the starting point and the transition point of a task session. As we discussed in Section II, various ways of starting a task, interrupting, and task switching a task make it a challenging problem. Coman's algorithm ignores interruption, therefore addresses only a simplified version of the sub-problem. Due to the incorrect assumption, the task sessions detected by the algorithm may be only stages of solving a task. For a simple task, a task may appear to be one stage as in the lab study. But in industry setting, a task often involves multiples stages, and how to link multiple stages is not clear.

The second sub-problem of linking task sessions remains unsolved, but we consider it closely related to linking multiple stages of a task. If we are able to link two stages through some criteria, such as common artifacts, we may also be able to link two sessions through the same criteria.

2) How to make it better?: Coman's algorithm may be improved by taking interruptions into consideration. One technique may be making the algorithm be aware of regular interruption time period, such as lunch time. When the DOA computation goes across the regular interruption period, it is handled differently. Another method is to "shrink" the interruption to reduce the gap effect, as we did in the case study. When an interruption occurs, the DOAs of all the methods decrease greatly and may result a gap in the TIIA time series that affects the task boundary detection. So if the interruption time periods are artificially shortened, say from longer than five minutes to one minute, then the decrease of DOA may not be large enough to end all the TIIAs and result a gap. At the same time, since there is still an interruption time period of one minute, the DOA still exhibits a decrement that reflects the effect of the interruption.

There is yet no solution to the second sub-problem of linking task sessions. We think that the key to solving this sub-problem is to find the link between the task sessions. The link may be common artifacts (being viewed or changed), closely related artifacts (such as peer concepts), code clones, or others. Task sessions with a large number of common artifacts that are ever accessed may belong to the same task. It is also likely that task sessions with a small number of artifacts being changed may belong to the same task. In principle, linkage criteria can make use of all kinds of information available in the interaction history, such as artifact, time, action (such as editing, searching, debugging ). When task sessions are linked, the purpose of task sessions may be also identified. For example, if a task session with many searches and viewing is followed by a session with a lot of editing, and if the artifacts that are viewed most in the first is changed in the second, then it may be the case that the first task session was to find task relevant information, while the second was to perform the appropriate changes.

We can also improve the algorithm by including new information sources into analysis. A bug tracking system, such as Bugzilla, is commonly used in a development team to track all the reported issues, or tasks. It records various information about each issue, such as priority, time being reported, and time being solved. All of this information may be used together with interaction history to improve the detection of task boundaries. Another information source is the version control system, such as CVS, subversion, or git. A version control system records all the changes to a software plus meta-data about the changes, including author identity, date, and possibly a reference to a related bug report. In many development teams, it is common to commit changes right after the task is solved, and input task information (or the bug ID) when committing changes. Therefore, depending on the tools used and the known practices of the development team, the rough time of a task may be recorded in the version control system and could be used in identifying task boundaries.

# D. Summary

Our study has shown that Coman's task splitting algorithm tends to detect more tasks than self-reported. Detailed data examination suggests that the underlying assumption about task core is unreasonable and the effects of interruption are not considered. Further analysis of the general problem of task boundary detection shows that the problem consists of two sub-problems. Coman's algorithm addresses the first sub-problem but still need improvement, while the second sub-problem remains unsolved.

### VI. RELATED WORK

# A. Mining Interaction history

Interaction history contains rich information about how program artifacts are accessed during software maintenance. Recent research suggests that it is a promising source for better understanding software development. For example, Schneider et al. proposed to study programmer interactions to improve team coordination in distributed projects [12]. Parnin et al. developed a technique to extract usage context by mining interaction histories [13]. Zou and Godfrey used interaction histories to detect couplings [14].

## B. Task-aware Software Development

Tasks are the work units of software development. Researchers have proposed to support tasks explicitly to assist programmers work. Murphy et al. suggested to improve IDEs and collaboration based on task structure — the subset of program artifacts and relations that are relevant to a task [3]. Kersten et al. used task context to recommend task relevant artifacts [15]. Robillard et al. proposed to support task-aware software development environment based on navigation analysis [6]. Röthlisberger et al. used heatmap to highlight artifacts within an IDE in a configurable way [16].

## VII. THREATS TO VALIDITY

In our study, the information of tasks are collected from programmers' self-reports at the end of the day. This assumes that programmers know what a task is, they can recall what they did, and they are willing to report it to us. This method may cause programmers to give a compressed report naturally, or only report their main tasks. We ignore non-responses without further asking why. This may cause data bias. We plan to address these issues in the future by performing complimentary studies.

As with all the case studies, this research has external validity threat. We chose this company only because it was available to us. The programmers participated the study because they were willing to do so. More case studies should be performed to answer our research questions in a more generalized basis.

### VIII. CONCLUSIONS AND FUTURE WORK

In this paper, we performed an industrial case study of Coman's automatic algorithm for detecting task boundaries from interaction histories. Based on the data from six programmers working for 12 days, we found that the algorithm performed quite differently comparing with the original lab validation study: many more tasks than self-reported were detected, and the best results after parameter tuning still had an error rate about 76%. Further analysis of the data led us discover two problems within Coman's algorithm: first, the underlying assumption about task core may not be reasonable in an industrial setting; and second, the demonstrable e?ects of interruptions are ignored. We discussed the general problem of task boundary detection and argued that the problem consists of two sub-problems: identifying task sessions and linking task sessions. Coman's algorithm in fact addresses the first sub-problem but under incorrect assumption. The second subproblem remains unsolved. We discuss possible techniques that may be used to improve Coman's algorithm and to address the second sub-problem.

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